

MAJORITY DYNAMICS AND THE RETENTION OF INFORMATION

OMER TAMUZ AND RAN J. TESSLER

ABSTRACT. We consider a group of agents connected by a social network who participate in *majority dynamics*: each agent starts with an opinion in $\{-1, +1\}$ and repeatedly updates it to match the opinion of the majority of its neighbors.

We assume that one of $\{-1, +1\}$ is the “correct” opinion S , and consider a setting in which the initial opinions are independent conditioned on S , and biased towards it. They hence contain enough information to reconstruct S with high probability. We ask whether it is still possible to reconstruct S from the agents’ opinions after many rounds of updates.

While this is not the case in general, we show that indeed, for a large family of bounded degree graphs, information on S is retained by the process of majority dynamics.

Our proof technique yields novel combinatorial results on majority dynamics on both finite and infinite graphs, with applications to zero temperature Ising models.

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1. INTRODUCTION

Consider a group of people (agents) who each carry one of two possible opinions regarding some issue. Each agent forms an initial opinion, and then repeatedly updates it by observing the opinions of its neighbors in a social network.

Many variants of this model have been studied in diverse settings and disciplines such as Economics (e.g. [1, 5]), Statistical Mechanics (e.g. [4, 6, 10]), Computer Science (see a survey by Shah [18]), and Mathematics (e.g. [3, 11]). These include variations on how the agents acquire their initial opinions (e.g., deterministically, at random, arbitrarily or through some other process), what they aim to achieve in this process (e.g., rational agents in Economics [1, 7, 16], message passing agents in Computer Science [12]), and how they go about updating their opinions in order to achieve this.

A particularly well studied model is that of *majority dynamics*, in which agents update their opinions to match that of the majority of their neighbors. We choose a setting in which each agent's initial opinion carries some independent information regarding the "true" opinion, and where it is possible to discover this truth with high probability by aggregating the initial opinions. The question we tackle is the following: is information lost through the process of majority dynamics? When can the "truth" be well estimated even after people have exchanged opinions?

These questions were, to our knowledge, first considered in the context of majority dynamics in Mossel, Neeman and Tamuz [15], who gave both positive and negative examples. We extend their work in several directions, as described below.

As part of our analysis we develop a combinatorial tool that is instrumental in studying majority dynamics. Using it, we present a number of purely combinatorial results on this process. These results can be applied directly to what is known as the dynamics of zero temperature translation-invariant ferromagnetic Ising model, which, on odd degree graphs, is the same as what we below call asynchronous majority dynamics. In particular, Theorem 1.5 gives, for every odd degree lattice of (say) polynomial growth, a constant C such that regardless of the starting configuration, with probability one no site changes its state more than C times.

1.1. Definitions and a statement of the problem. Let V be a finite or countably infinite set of agents. Let $G = (V, E)$, the *social network*, be an undirected, connected, locally finite graph. We denote the neighbors of $i \in V$ by $N(i) = \{j : (i, j) \in E\}$, and say that G is d -bounded degree when $|N(i)| \leq d$ for all $i \in V$.

We denote by $A_t^i \in \{-1, +1\}$ agent i 's *opinion* at time t . After drawing some initial opinions $\{A_0^i\}_{i \in V}$ from a distribution we describe below, the agents proceed to update their opinions using *majority dynamics*.

We consider two version of majority dynamics, namely a synchronous and an asynchronous one, and prove all our results for both. In the synchronous model, t takes values in the non-negative integers only, and we set

$$A_{t+1}^i = \operatorname{sgn} \sum_{j \in N(i)} A_t^j. \quad (1.1)$$

In the asynchronous model, we equip each agent with an independent, unit rate Poisson clock, and let each agent update its opinions at the times of its clock rings,

using

$$A_t^i = \operatorname{sgn} \sum_{j \in N(i)} A_t^j. \quad (1.2)$$

We assume throughout that the degree of every vertex is odd, so that there are no “ties” and $A_t^i \in \{-1, +1\}$ for all t . Alternatively, given a graph with even degrees, one could add i or remove i from $N(i)$ to make $|N(i)|$ odd.

In the asynchronous model, in order to ensure that the model is well defined, we must rule out the possibility that there will occur an infinite sequence of clock ticks $t_1 > t_2 > \dots > 0$ such that agent i updates at time t_i and $i \in N(i+1)$. To ensure that this occurs with probability zero, it is sufficient to assume that degrees are bounded. We indeed make this assumption throughout, and prove that it is sufficient in Claim 3.5 below.

A classical result [13] is that for finite graphs, in the synchronous model, each agent’s opinion either converges, or, from some time on, oscillates between -1 and $+1$ with period two; that is, each agent’s opinion eventually has period at most two. In the asynchronous model all opinions converge for finite graphs.

The same can be shown to hold for bounded degree infinite graphs with sufficiently slow growth. To state this result we shall need some definitions. Denote graph distances by $\rho(\cdot, \cdot)$, let $n_r(G, i) = |\{j : \rho(i, j) = r\}|$ be number of vertices at graph distance r from i in G . Finally, denote

$$M_d[G, i] = \sum_{r=0}^{\infty} \left(\frac{d+1}{d-1} \right)^{-r} n_r(G, i);$$

the usefulness of this definition will be demonstrated below.

Ginosar and Holzman [8] prove the following result, which is a strengthening of a somewhat weaker claim by Moran [14].

Theorem 1.1 (Moran, Ginosar and Holzman¹). *Let G be a d -bounded degree graph such that $M_d[G, i] < \infty$ for some (\Leftrightarrow all) $i \in V$. Then each agent’s opinion eventually has period at most two in the synchronous model, and converges in the asynchronous model.*

Note that this is a combinatorial (rather than a probabilistic) result, in the sense that it holds for every set of initial opinions $\{A_0^i\}_{i \in V}$, and, in the asynchronous model, for every choice of clock ring times (in which there are no rings at the same time). We prove this result using a novel combinatorial tool, which yields additional insights into this process (see Section 1.2.1). In particular, we bound the number of times that an agent may change its opinion, and show that an agent will never change its opinion if a large enough neighborhood around it agrees with it.

In Section 3.7 we provide an example of a d -regular graph for which $M_d[G, i] = \infty$ for all $i \in V$, and for which, in the synchronous model, each agent’s opinion eventually has period at most two for any initial set of opinions. This shows that the sufficient condition of Theorem 1.1 is not necessary. Indeed, the existence of a simple, geometrical, necessary and sufficient condition for convergence to period at most two is an interesting open question.

It should be noted that had we allowed even degrees with random tie breaking, the dynamics would have changed dramatically. Tessler and Louidor [19] discuss

¹Note that Moran, as well as Ginosar and Holzman, prove their theorems for the synchronous model; we extend it to the asynchronous model.

the asynchronous model for trees with even degrees larger than two (for which it may be that $M_d[G, i] < \infty$). It is shown there that if the initial opinions are chosen with i.i.d symmetric Bernoulli distribution, then almost surely there exist some agents that change their opinions infinitely many times.

We henceforth consider only *slow growth* graphs, i.e. graph for which

$$M_d[G, i] < \infty,$$

and for which, by Theorem 1.1, we can define random variables Z^i taking values in $\{-1, +1\}$ and given by

$$Z^i = \lim_{t \rightarrow \infty} A_{2t}^i,$$

where the limit is taken over $t \in \mathbb{N}$.

We next describe how the agents acquire their initial opinions. Let $S \in \{-1, +1\}$ be the *state of the world* with $\mathbb{P}[S = -1] = \mathbb{P}[S = 1] = \frac{1}{2}$. We think of $A_t^i \in \{-1, +1\}$ as agent i 's *opinion regarding* S at time t , and draw $\{A_0^i\}_{i \in V}$ as follows: We fix some $\frac{1}{2} < p < 1$, and let $A_0^i = S$ with probability p and $A_0^i = -S$ with probability $1 - p$, with the events ' $A_0^i = S$ ' being independent of each other and of S . Note that it follows that the random variables A_0^i are not independent, but are independent (in fact, i.i.d) conditioned on S .

When $|V|$ is finite but large, then S can be estimated with high probability given $\{A_0^i\}_{i \in V}$. When $|V|$ is infinite, S can be estimated exactly given $\{A_0^i\}_{i \in V}$. Formally, let $|V| = n$, and let

$$\hat{S}_0 = \operatorname{sgn} \sum_{i \in V} A_0^i.$$

Then by the Chernoff bound

$$\mathbb{P}[\hat{S}_0 \neq S] \leq e^{-O(n)}.$$

For the case that $|V| = \infty$, let

$$\mathcal{A}_0 = \sigma\left(\{A_0^i\}_{i \in V}\right).$$

Here we denote by $\sigma(X)$ the sigma-algebra generated by X , a random variable². Then it holds that

$$\inf_{\hat{S}_0 \in \mathcal{A}_0} \mathbb{P}[\hat{S}_0 \neq S] = 0,$$

where, by a slight abuse of notation, we say that $\hat{S}_0 \in \mathcal{A}_0$ when \hat{S}_0 is \mathcal{A}_0 -measurable.

The question that we tackle is the following: when is it the case that S can still be reconstructed from the limiting opinions $\{Z^i\}_{i \in V}$? Formally, define $\mathcal{A}_\infty = \sigma\left(\{Z^i\}_{i \in V}\right)$, and let the probability of error (in the reconstruction of S using $\{Z^i\}_{i \in V}$) be given by

$$\delta(G, p) = \inf_{\hat{S} \in \mathcal{A}_\infty} \mathbb{P}[\hat{S} \neq S],$$

where, to remind the reader, $p = \mathbb{P}[A_0^i = S]$ is a parameter of our measure \mathbb{P} .

We ask the question of whether $\delta(G, p)$ is equal to zero in the case of an infinite graph, or is close to zero in the case of a large finite graph. Formally, let $\{G_n\}_{n \in \mathbb{N}}$

²This is the smallest sigma-algebra for which X is measurable.

be a sequence of finite graphs such that $\lim_n |V_n| = \infty$. For which sequences and values of p is it the case that $\lim_n \delta(G_n, p) = 0$? And for which infinite G does it hold that $\delta(G, p) = 0$?

Berger [2] gives an example of a sequence of finite graphs with $\lim_n |V_n| = \infty$ and such that in each graph there exists a *dynamic monopoly* of size eighteen: a set $W \subset V$ of eighteen vertices with the property that if $\{A_0^j\}_{j \in W}$ are all equal to some $s \in \{-1, +1\}$ then $Z^i = s$ for all $i \in V$ (in the synchronous model)³. Since $\mathbb{P}[A_0^j = -S \ \forall j \in W] = (1-p)^{18}$, it follows that $\delta(G_n, p) \geq (1-p)^{18}$.

This example involves graphs with increasingly large degrees. We offer the following conjecture (see also [15]).

Conjecture 1.2. *Let $\{G_n\}_{n \in \mathbb{N}}$ be a sequence of d -bounded degree finite graphs such that $\lim_n |V_n| = \infty$. Then*

$$\lim_{n \rightarrow \infty} \delta(G_n, p) = 0$$

for all $\frac{1}{2} < p < 1$.

Mossel, Neeman and Tamuz [15] show that if $\{G_n\}$ is sequence of *transitive* graphs (that is, graphs in which all nodes play the same role, see Section 1.2.2 below) then $\lim_n \delta(G_n, p) = 0$. In fact, they show that the same holds under a weaker assumption, namely that each geometric equivalence class is large (again, see Section 1.2.2 for a precise definition). They also show the same for good enough *expander* graphs.

For infinite graphs G , our question is whether $\delta(G, p) = 0$. While we provide some positive results below, we do not know the answer to the following seemingly basic question:

Question 1.3. *Does there exist an infinite graph G and $\frac{1}{2} < p < 1$ such that $\delta(G, p) > 0$?*

We prove the following claim⁴ (see [17] for related work).

Theorem 1.4. *In the synchronous model, no infinite, locally finite graph has a dynamic monopoly of finite size.*

This result suggests that perhaps $\delta(G, p) = 0$ for every finite degree, slow growth graph G .

1.2. Results.

1.2.1. *Majority Dynamics.* We begin by proving two combinatorial claims regarding majority dynamics. These claims may be of independent interest, but are also useful in proving our main results.

The first result is a quantitative version of Theorem 1.1.

Theorem 1.5. *Let G be a d -bounded degree graph such that $M_d[G, i] < \infty$ for some (\leftrightarrow all) $i \in V$. Then*

- (1) *In the synchronous model, the number of times t for which $A_{t+1}^i \neq A_{t-1}^i$ is at most $\frac{d+1}{d-1} \cdot d \cdot M_d[G, i]$.*

³We slightly weaken his definition; he requires that $A_t^i = s$ for all i in some finite time t .

⁴Although the proof of this claim is rather straightforward, we have not found it in the literature.

(2) In the asynchronous model, the number of times in which i changes its opinion is at most $\frac{d+1}{d-1} \cdot 2d \cdot M_d[G, i]$.

We would like to thank the anonymous referee for pointing out to us that this theorem can, for the synchronous model, be easily derived from the work of Ginosar and Holzman.

The next result shows that if a sufficiently large neighborhood of i starts with a certain opinion then i will always have this opinion.

Theorem 1.6. *Let G be a d -bounded degree graph such that $M_d[G, i] < \infty$ for some (\leftrightarrow all) $i \in V$. Let r_0 be such that*

$$\frac{d+1}{d-1} \cdot 2d \cdot \sum_{r>r_0} \left(\frac{d+1}{d-1}\right)^{-r} n_r(G, i) < 1.$$

If $A_T^j = A_T^i$ for some T and for all j such that $\rho(i, j) \leq r_0 + 2$, then $A_t^i = A_T^i$ for all $t > T$.

1.2.2. *Invariant random subgraphs.* A graph automorphism of $G = (V, E)$ is a bijection $h : V \rightarrow V$ such that $(i, j) \in E \leftrightarrow (h(i), h(j)) \in E$. We denote by $\text{Aut}(G)$ the automorphism group of G . Indeed, it is easy to verify that the graph automorphisms of a given graph form a group under composition.

Let H be a subgroup of $\text{Aut}(G)$. We say that H acts *transitively* on (the vertices of) G when, for every $i, j \in V$ there exists an $h \in H$ such that $h(i) = j$. Equivalently, H acts transitively on G when V/H , the set of H orbits of V , is a singleton. We say that H acts *quasi-transitively* on G when V/H is finite. Finally, we say that G is (quasi-) transitive when $\text{Aut}(G)$ acts on it (quasi-) transitively.

A subgraph of $G_0 = (V_0, E_0)$ is a graph $G = (V, E)$ such that $V \subseteq V_0$ and $E \subseteq E_0$ is a set of edges on V . Let H be a subgroup of $\text{Aut}(G_0)$. A random G_0 -subgraph G (that is, a random variable G that takes values in the space of subgraphs of a graph G_0) is said to have an *H -invariant* distribution if for all $h \in H$ the law of $h(G)$ equals the law of G .

Note that G could have nodes with even degrees even when G_0 has odd degrees. Since we will want to apply majority dynamics to G , we add or remove self-loops to G in order to make all degrees odd. This does not affect the fact that G is H -invariant; the modified G is H -invariant iff the unmodified G was H -invariant. This also does not increase any degree beyond what it was on G_0 , since we do not add a self-loop unless we remove another edge.

The following is our main result for this section.

Theorem 1.7. *Let G_0 be a quasi-transitive infinite graph with maximal degree d such that $M_d[G_0, i] < \infty$ for some (\leftrightarrow all) $i \in V_0$.*

Let $H \leq \text{Aut}(G_0)$ act quasi-transitively on G_0 , and let G be an infinite connected random subgraph of G_0 with an H -invariant distribution. Then

$$\delta(G, p) = 0$$

almost surely for any $\frac{1}{2} < p < 1$.

It follows from Theorem 1.7 that $\delta(G, p) = 0$ for any quasi-transitive, slow growth infinite graph G and any $\frac{1}{2} < p < 1$. This is already a non-trivial result, which is generalized in Theorem 1.7.

1.2.3. *Families of uniformly bounded growth graphs.* Given $d \geq 3$ and a function $f : \mathbb{N} \rightarrow \mathbb{N}$, let $\mathcal{B}(f, d)$ be the family of d -bounded degree graphs G such that $n_r(G, i) \leq f(r)$ for all vertices i in G and $r \in \mathbb{N}$.

Reusing the notation $M_d[\cdot]$, let

$$M_d[f] = \sum_{r=0}^{\infty} \left(\frac{d+1}{d-1} \right)^{-r} \cdot f(r).$$

We say that $f : \mathbb{N} \rightarrow \mathbb{N}$ has *slow growth* if $M_d[f] < \infty$.

Theorem 1.8. *Fix an odd $d \geq 3$, and let $f : \mathbb{N} \rightarrow \mathbb{N}$ have slow growth. There exists a $\frac{1}{2} < p_0 < 1$ such that for all $p_0 < p < 1$ and for all sequences of finite graphs $\{G_n\}_{n \in \mathbb{N}}$ in $\mathcal{B}(f, d)$ such that $\lim_n |V_n| = \infty$ it holds that*

$$\lim_n \delta(G_n, p) = 0,$$

and for all infinite graphs $G \in \mathcal{B}(f, d)$ it holds that

$$\delta(G, p) = 0.$$

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3. PROOFS

3.1. **The Lyapunov functional.** Following Goles and Olivos [9], we analyze the process of Majority dynamics using the technique of Lyapunov functionals. In particular, we build and elaborate on the ideas of Ginosar and Holzman [8] to apply these techniques to infinite graphs. We apply them to both the asynchronous and the synchronous model, and use them to prove some additional, general results for majority dynamics on graphs with uniformly bounded growth.

We begin by introducing some non-standard notation which will allow us to simultaneously treat the synchronous and the asynchronous models. We let the symbol Δ stand for “1” for in synchronous model and for “ dt ” in the asynchronous model. For the asynchronous model, we denote

$$A_{t-\Delta}^i = A_{t-dt}^i = \lim_{t' \rightarrow t^-} A_{t'}^i$$

and

$$A_{t+\Delta}^i = A_{t+dt}^i = \lim_{t' \rightarrow t^+} A_{t'}^i = A_t^i,$$

where the last equality follows from Eq. 1.2. The definitions of majority dynamics, Eqs. 1.1 and 1.2, can now be written in one equation:

$$A_{t+\Delta}^i = \text{sgn} \sum_{j \in N(i)} A_t^j.$$

Let $G = (V, E)$ be a finite or infinite d -bounded degree graph. Let a d -legal edge weighting $z : E \rightarrow (0, 1]$ be such that for any two adjacent edges $e_1 = (i, j)$ and $e_2 = (i, k)$ it holds that

$$\frac{z(e_1)}{z(e_2)} < 1 + \frac{2}{d-1} = \frac{d+1}{d-1}. \quad (3.1)$$

Note that z is a function on the undirected edge set E , and so $z(i, j) = z(j, i)$. It is easy to see that

$$\operatorname{sgn} \sum_{j \in N(i)} A_t^j = \operatorname{sgn} \sum_{j \in N(i)} z(i, j) A_t^j. \quad (3.2)$$

Indeed, assume i has k neighbors, and that the majority of their opinions are, without loss of generality, $+1$. Then i has at least $\frac{k+1}{2}$ neighbors with opinion $+1$. Denote by N_+ this set of neighbors. Similarly, i has at most $\frac{k-1}{2}$ neighbors with opinion -1 , which we denote by N_- . Let z_1 be the largest $z(i, j)$ for $j \in N(i)$, and z_2 be the smallest $z(i, j)$ in the same set. By assumption, and because $k \leq d$,

$$\frac{z_1}{z_2} < \frac{d+1}{d-1} \leq \frac{k+1}{k-1}.$$

Hence,

$$\begin{aligned} \sum_{j \in N(i)} z(i, j) A_t^j &\geq z_2 \sum_{j \in N_+} 1 - z_1 \sum_{j \in N_-} 1 \\ &= z_2 \left(|N_+| - \frac{z_1}{z_2} |N_-| \right) \\ &\geq z_2 \left(\frac{k+1}{2} - \frac{z_1}{z_2} \frac{k-1}{2} \right), \end{aligned}$$

the last expression is positive due to the second inequality in the previous display, and Eq. 3.2 follows. Thus, the definitions of majority dynamics (Eqs. 1.1 and 1.2) can equivalently be written as

$$A_{t+\Delta}^i = \operatorname{sgn} \sum_{j \in N(i)} z(i, j) A_t^j.$$

Yet another equivalent definition is

$$A_{t+\Delta}^i = \operatorname{argmin}_{a \in \{-1, +1\}} \sum_{j \in N(i)} z(i, j) (A_t^j - a)^2; \quad (3.3)$$

intuitively, each agent in each turn can be seen as trying to minimize the “energy” $\sum_{j \in N(i)} z(i, j) (A_t^j - a)^2$. This motivates the following definition:

$$L_t = \frac{1}{4} \sum_{(i, j) \in E} z(i, j) (A_{t+\Delta}^i - A_t^j)^2.$$

Of course, L_t has to be finite for this definition to be useful, and we indeed give in Proposition 3.4 a necessary and sufficient geometrical condition under which a d -legal summable z exists. Note also that L_t depends (implicitly) on the choice of z .

It turns out that this definition of L_t is a correct choice for a “Lyapunov functional”, in the sense that L_t is monotone non-increasing.

Proposition 3.1. $L_{T+t} \leq L_T$ for all $T, t \geq 0$.

Before proving this proposition, we will need the following definitions and lemma. Define J_t^i by

$$J_t^i = \frac{1}{2} (A_{t+\Delta}^i - A_{t-\Delta}^i) \sum_{j \in N(i)} z(i, j) A_t^j,$$

and let

$$J_t = \sum_{i \in V} J_t^i.$$

Claim 3.2. $J_t^i \geq 0$, and $J_t^i = 0$ iff $A_{t+\Delta}^i = A_{t-\Delta}^i$.

Proof. Since we assume that $|N(i)|$ is odd, $\sum_{j \in N(i)} z(i, j) A_t^j$ is never zero. It follows that $J_t^i = 0$ iff $A_{t+\Delta}^i = A_{t-\Delta}^i$.

To see that $J_t^i \geq 0$, note that when $A_{t+\Delta}^i = A_{t-\Delta}^i$ then $J_t^i = 0$. Otherwise we have that $J_t^i = A_{t+\Delta}^i \sum_{j \in N(i)} z(i, j) A_t^j$. But $A_{t+\Delta}^i = \text{sgn} \sum_{j \in N(i)} z(i, j) A_t^j$ and so J_t^i is equal to the product of two equal sign multiplicands and is therefore positive. \square

Claim 3.3.

$$L_t - L_{t-\Delta} = -J_t.$$

Proof.

$$\begin{aligned} L_t - L_{t-\Delta} &= \frac{1}{4} \sum_{(i,j) \in E} z(i, j) \left[(A_{t+\Delta}^i - A_t^j)^2 - (A_t^i - A_{t-\Delta}^j)^2 \right] \\ &= -\frac{1}{2} \sum_{(i,j) \in E} z(i, j) \left[A_{t+\Delta}^i A_t^j - A_t^i A_{t-\Delta}^j \right] \end{aligned}$$

Since the edges are undirected, summing over $(i, j) \in E$ is the same as summing over $(j, i) \in E$. Therefore, and since $z(i, j) = z(j, i)$, we can exchange the roles of i and j in the last summand:

$$\begin{aligned} &= -\frac{1}{2} \sum_{(i,j) \in E} z(i, j) \left[A_{t+\Delta}^i A_t^j - A_t^j A_{t-\Delta}^i \right] \\ &= -\sum_{i \in V} \frac{1}{2} (A_{t+\Delta}^i - A_{t-\Delta}^i) \sum_{j \in N(i)} z(i, j) A_t^j \\ &= -\sum_{i \in V} J_t^i \\ &= -J_t. \end{aligned}$$

\square

The proof of Proposition 3.1 is now immediate.

Proof of Proposition 3.1. Since $L_t - L_{t-\Delta} = -J_t$ and since $J_t \geq 0$, it follows that $L_{T+t} \leq L_T$ for all $T, t \geq 0$. \square

3.2. No dynamic monopolies on infinite graphs. As a simple application of Claims 3.2 and 3.3 we show that infinite graphs cannot have dynamic monopolies.

Proof of Theorem 1.4. Consider the synchronous model. Let G be an infinite graph, and let $z : E \rightarrow \mathbb{R}^+$ be the d -legal edge weighting given by the constant function $z = 1$. Then z is not summable, but Claims 3.2 and 3.3 still hold, given that initial signals are chosen so that L_t is finite.

Let W be a finite set of vertices. Let $A_0^i = +1$ for all $i \in W$ and $A_0^i = -1$ for all $i \notin W$. Then L_0 is finite. Note that with z constant, L_t and J_t^i are integer, and therefore L_t decreases by at least one whenever $A_{t+1}^i \neq A_{t-1}^i$. Since $L_t \geq 0$, at

most a finite number of agents change their opinion to +1, and W is not a dynamic monopoly. \square

It may be possible to give a stronger, quantitative version of this theorem, by using an edge weighting that increases with the distance from W .

3.3. Summable edge weightings and slow growth. We are now almost ready to prove Theorem 1.1. Before that, we will show that graphs with slow growth admit summable, d -legal edge weightings.

Proposition 3.4. *Let $G = (V, E)$ be a d -bounded degree graph. Then the following are equivalent.*

- (1) G admits a summable, d -legal edge weighting.
- (2) $M_d[G, i] < \infty$ for some (\leftrightarrow all) $i \in V$.

Proof. Let G be a d -bounded degree graph, and denote $a = (d + 1)/(d - 1)$. For a node i and an edge $e = (j, k)$, denote by $\rho(i, e) = \min\{\rho(i, j), \rho(i, k)\}$ the distance of e from i .

We first show that (1) implies (2). Let G admit a summable, d -legal edge weighting z . Let $e_0 = (i, j)$ an edge. By the definition of d -legal edge weightings (Eq. 3.1), given an edge e such that $\rho(i, e) = r$, we have that

$$z(e) \geq z(e_0)a^{-r-1}.$$

Hence, by the fact that z is summable, it follows that

$$\infty > \sum_{e \in E} z(e) = \sum_{r=0}^{\infty} \sum_{\{e \in E : \rho(i, e) = r\}} z(e) \geq \frac{z(e_0)}{a} \sum_{r=0}^{\infty} a^{-r} |\{e \in E : \rho(i, e) = r\}|.$$

Now, the number of edges at distance r is greater than or equal to the number of nodes at distance $r + 1$, $n_{r+1}(G, i)$. Hence

$$\begin{aligned} \infty &> \frac{z(e_0)}{a} \sum_{r=0}^{\infty} n_{r+1}(G, i) a^{-r} \\ &= z(e_0) \sum_{r=1}^{\infty} n_r(G, i) a^{-r} \\ &= z(e_0) (M_d[G, i] - 1), \end{aligned}$$

where the last equality follows from the definition of $M_d[G, i]$ and the fact that $n_0(G, i) = 1$. Thus $M_d[G, i]$ is finite.

We now show that (2) implies (1). Let $M_d[G, i] < \infty$ for some $i \in V$. Let $\eta : \mathbb{N} \rightarrow (1/a, 1)$ be any monotone increasing function, and let

$$z(e) = a^{-\rho(i, e)} \frac{\eta(\rho(i, e))}{\eta(0)}$$

be an edge weighting. We will show that it is d -legal and summable. Indeed,

$$\begin{aligned} \sum_{e \in E} z(e) &= \sum_{r=0}^{\infty} a^{-r} \frac{\eta(r)}{\eta(0)} |\{e \in E : \rho(i, e) = r\}| \\ &\leq \sum_{r=0}^{\infty} a^{-r} \frac{\eta(r)}{\eta(0)} \cdot d \cdot |\{j \in V : \rho(i, j) = r\}|, \end{aligned}$$

since the number of edges at distance r is at most d times the number of vertices at that distance. Since $n_r(G, i) = |\{j \in V : \rho(i, j) = r\}|$ and since $\eta(r) < 1$ then

$$\begin{aligned} &< \frac{d}{\eta(0)} \sum_{r=0}^{\infty} a^{-r} \cdot n_r(G, i) \\ &= \frac{d}{\eta(0)} \cdot M_d[G, i] \\ &< \infty, \end{aligned}$$

and so z is summable.

To see that z is d -legal, note that if e_1 and e_2 are adjacent then either

$$\rho(i, e_1) = \rho(i, e_2)$$

in which case $z(e_1) = z(e_2)$, or else, without loss of generality

$$\rho(i, e_2) = \rho(i, e_1) + 1.$$

In this case, denoting $r = \rho(i, e_1)$,

$$\frac{z(e_2)}{z(e_1)} = \frac{a^{-r-1}\eta(r+1)}{a^{-r}\eta(r)} = a^{-1} \frac{\eta(r+1)}{\eta(r)} < 1,$$

where the last inequality follows from the fact that $\eta(r+1) < 1$ and $\eta(r) > 1/a$. Likewise,

$$\frac{z(e_1)}{z(e_2)} = \frac{a^{-r}\eta(r)}{a^{-r-1}\eta(r+1)} = a \frac{\eta(r)}{\eta(r+1)} < a = \frac{d+1}{d-1},$$

where the last inequality follows from the fact that η is monotone increasing. \square

Proof of Theorem 1.1. Let z be a d -legal summable edge weighting, as defined in the proof of Proposition 3.4. Note that L_t is non-negative, and also finite for all t :

$$\begin{aligned} L_t &= \frac{1}{4} \sum_{(i,j) \in E} z(i, j) (A_{t+\Delta}^i - A_t^j)^2 \\ &\leq \sum_{(i,j) \in E} z(i, j). \end{aligned}$$

Since L_t is finite, non-negative and non-increasing, it follows that

$$L = \lim_{t \rightarrow \infty} L_t$$

always exists and is non-negative.

Fix the initial opinions, the times of the clock rings (for the asynchronous model) and a vertex j . Note that J_t^j is either zero or greater than some $\epsilon_j > 0$. Let T be such that $L_T - L < \epsilon_j$. It then follows from Claims 3.2 and 3.3 that $J_t^j = 0$ for all $t > T$, since otherwise it would be the case that $J_t^j > \epsilon_j$ and $L_{t+\Delta} < L$. Now, by Claim 3.2 it follows from this that $A_{t+\Delta}^j = A_{t-\Delta}^j$ for all $t > T$, and so j 's opinion has period at most two in the synchronous case, and converges in the asynchronous case. \square

3.4. Combinatorial majority dynamics results. Let $G = (V, E)$ be a d -bounded degree graph such that $M_d[G, i] < \infty$ for some $i \in V$. As per the proof of Proposition 3.4, let z be a summable, d -legal edge weighting given by

$$z(e) = a^{-\rho(i,e)} \frac{\eta(\rho(i,e))}{\eta(0)},$$

where $a = (d+1)/(d-1)$, and $\eta : \mathbb{N} \rightarrow (1/a, 1)$ is some monotone increasing function. In fact, by the proof of Proposition 3.4, we have that

$$\sum_{e \in E} z(e) \leq \frac{d}{\eta(0)} \cdot M_d[G, i] \leq \frac{d+1}{d-1} \cdot d \cdot M_d[G, i].$$

Proof of Theorem 1.5. By Claim 3.2, the number of times that $A_{t+\Delta}^i \neq A_{t-\Delta}^i$ is equal to the number of times that $J_t^i \neq 0$. Hence we will prove the claim by bounding the number of times that $J_t^i \neq 0$.

Since $z(i, j) = 1$ for all $j \in N(i)$, $J_t^i \geq 1$ whenever $J_t^i \neq 0$. Since $L = \lim_{t \rightarrow \infty} L_t$ is non-negative, and since $L_t - L_{t-\Delta} = -J_t$ by Claim 3.3, it follows that the number of times that $J_t^i \neq 0$ is at most L_0 .

Now, in the asynchronous case

$$\begin{aligned} L_0 &= \frac{1}{4} \sum_{(k,j) \in E} z(k,j) (A_0^k - A_0^j)^2 \\ &\leq 2 \sum_{e \in E} z(e) \\ &\leq \frac{d+1}{d-1} \cdot 2d \cdot M_d[G, i]. \end{aligned}$$

In the synchronous case this bound can be improved. By Eq. 3.3, we have that

$$\frac{1}{4} \sum_{j \in N(k)} z(j,k) (A_1^k - A_0^j)^2 < \frac{1}{2} \sum_{j \in N(k)} z(j,k).$$

Hence

$$\begin{aligned} L_0 &= \frac{1}{4} \sum_{(k,j) \in E} z(k,j) (A_1^k - A_0^j)^2 \\ &< \sum_{e \in E} z(e) \\ &\leq \frac{d+1}{d-1} \cdot d \cdot M_d[G, i]. \end{aligned}$$

□

Proof of Theorem 1.6. Let r_0 be such that

$$\frac{d+1}{d-1} \cdot 2d \cdot \sum_{r > r_0} \left(\frac{d-1}{d+1} \right)^r \cdot n_r(G, i) < 1.$$

Then it is easy to verify that

$$\sum_{e \in E} \mathbb{1}_{\rho(i,e) > r_0} z(e) < \frac{1}{2}.$$

If $A_T^j = s$ for all j within distance $r_0 + 2$ from i , then $A_{T+\Delta}^j = s$ for all j within distance $r_0 + 1$ from i , and $A_{T+\Delta}^k = s$ for all $e = (j, k)$ with $\rho(i, e) \leq r_0$. Hence

$$\begin{aligned} L_T &= \frac{1}{4} \sum_{(k,j) \in E} z(k, j) (A_{T+\Delta}^k - A_T^j)^2 \\ &\leq 2 \sum_{e \in E} \mathbb{1}_{\rho(i, e) > r_0} z(e) \\ &< 1. \end{aligned}$$

Now, as in the proof of Theorem 1.5 above, if $J_t^i \neq 0$ then $J_t^i \geq 1$. However, since $L_T < 1$, since $L_t - L_{t-\Delta} = -J_t$, and since L_t is non-negative, it follows that $J_t^i = 0$ for all $t > T$, and so, in particular, $A_{t+\Delta}^i = A_{t-\Delta}^i$ for all $t > T$. This completes the proof for the asynchronous model. In the synchronous model, $A_T^i = A_{T+1}^i$, since all of i 's neighbors also have $A_T^j = s$, and so it follows that $A_T^i = s$ for all $t > T$. \square

3.5. Invariant random subgraphs.

3.5.1. Light cones. We commence by defining for each vertex i and time t the set of vertices C_t^i that form i 's *past light cone* at time t . This is the set of vertices whose initial opinions may have influenced i 's opinion at time t . Formally, set $C_0^i = \{i\}$ for all $i \in V$. At every time t in which i updates its opinion, update C_t^i by

$$C_t^i = \bigcup_{j \in N(i)} C_{t-\Delta}^j.$$

In the synchronous model, C_t^i is simply the ball of radius t around i . In the asynchronous model every vertex is a member of C_t^i with positive probability. However C_t^i is still finite with probability one.

Claim 3.5. *Let G be a d -bounded degree graph. Then for all $i \in V$ and $t \geq 0$ it holds that $\mathbb{P}[|C_t^i| < \infty] = 1$.*

It follows that the asynchronous model is well defined, since if there exists an infinite sequence of times $t_1 > t_2 > \dots > 0$ such that i updates at time t_i and $i \in N(i+1)$, then $|C_{t_1}^i| = \infty$.

Proof. Note that $|C_t^i|$ is stochastically dominated from above by the total number of offsprings in a Galton-Watson process, where each vertex has d children with probability equal to the probability of a clock ring in $[0, t]$, and zero children otherwise. Since, for t small enough, the expected number of children in this process is lower than one, it follows that the total number of offsprings is almost surely finite. Hence for t small enough (e.g., $1/(100d)$) we have that $\mathbb{P}[|C_t^i| = \infty] = 0$.

Now, the event that $|C_t^i| = \infty$ is equivalent to the existence infinite sequence of times $t_1 > t_2 > \dots > 0$ such that (after an appropriate renaming of the vertices) i updates at time t_i and $i \in N(i+1)$. By the time-shift-invariance of the clock tick process we can assume that $\lim_i t_i = 0$, and so, if $\mathbb{P}[|C_t^i| = \infty] > 0$ for some t then $\mathbb{P}[|C_{1/(100d)}^j| = \infty] > 0$ for some j . But this is false by the above, and so $\mathbb{P}[|C_t^i| = \infty] = 0$ for all i and all t . \square

We say that vertices i and j are *causally connected* at time t if $C_t^i \cap C_t^j \neq \emptyset$. The following is immediate.

Claim 3.6. *Condition on the event that the clock ring times are such that i and j are not causally connected at time t . Then the events ‘ $A_t^i = S$ ’ and ‘ $A_t^j = S$ ’ are independent.*

3.5.2. *Choosing independent vertices.* We next proceed to construct, for each time t , a set of vertices whose opinions at time t are close to being independent, conditioned on S .

Let $G = (V, E)$ be an infinite quasi-transitive graph. Then it is standard to show that given $\delta > 0$ and a positive integer t , one may find a number $r_{t,\delta}$ such that, if i, j are two vertices whose distance is more than $r_{t,\delta}$, then the probability that i and j are causally connected at time t (meaning $C_t^i \cap C_t^j \neq \emptyset$) is at most δ . Note that in the synchronous case one can take $r_{t,\delta} = r_{t,0} = 2t$.

For every time t and $\delta > 0$, let $W_{t,\delta} \subseteq V$ be a random subset of V , drawn as follows. Associate to each vertex in $i \in V$ an i.i.d. exponential random variable X_i . A vertex i belongs to $W_{t,\delta}$ if and only if $X_i < X_j$ for all j with $\rho(i, j) \leq r_{t,\delta}$.

We note a number of easily verifiable facts regarding $W_{t,\delta}$.

- (1) Any two elements of $W_{t,\delta}$ are at distance more than $r_{t,\delta}$. Hence, by the definition of causal connectedness, any two vertices in $W_{t,\delta}$ are causally connected with probability at most δ .
- (2) For all $h \in \text{Aut}(G)$ it holds that $h(W_{t,\delta})$ has the same law as $W_{t,\delta}$; that is, the distribution of $W_{t,\delta}$ is $\text{Aut}(G)$ -invariant.
- (3) $W_{t,\delta}$ is almost surely infinite, since in quasi-transitive graphs the size of radius r balls is uniformly bounded.

If furthermore $H \leq \text{Aut}(G)$ acts quasi-transitively on G then $W_{t,\delta}$ intersects each orbit V/H infinitely often.

3.5.3. *Convergence and uniform convergence to Z .* Let $G = (V, E)$ be a slow-growth graph. Then, as we show above, the limit $Z^i = \lim_t A_{2t}^i$ exists almost surely. Hence for each vertex i the function q_i given by $q_i(2t) = \mathbb{P}[Z^i \neq A_{2t}^i]$ converges to zero.

Clearly, if i and j are both in the same orbit $V/\text{Aut}(G)$, then $q_i = q_j$. It follows that for every quasi-transitive graph, that is, for every graph with a finite number of orbits, there exists a single function q that converges to zero and such that $q(2t) \geq \mathbb{P}[Z^i \neq A_{2t}^i]$ for all $i \in V$. In this case we say that we have uniform convergence to Z^i . The next claim states that the same holds on a shift-invariant random graph chosen from a quasi-transitive graph.

Claim 3.7. *Let G_0 be a quasi-transitive infinite graph with maximal degree d such that $M_d[G_0, i] < \infty$ for some (\leftrightarrow all) $i \in V_0$.*

Let $H \leq \text{Aut}(G_0)$ act quasi-transitively on G_0 , and let G be an infinite random subgraph of G_0 with an H -invariant distribution. Then there exists a function $q : 2\mathbb{N} \rightarrow [0, 1]$ with $\lim_n q(2n) = 0$ such that $q(2t) \geq \mathbb{P}[Z^i \neq A_{2t}^i | i \in G]$.

Proof. Let $i, j \in V_0$ belong to the same orbit in V/H , so that there exists an $h \in H$ such that $h(i) = j$.

Since G is h -invariant, it is possible, using h , to couple two copies of our probability space in such a way that A_t^i in the first copy equals A_t^j in the second for all t , Z^i in the first copy equals Z^j in the second, and furthermore $i \in G$ in the first copy iff $j \in G$ in the second. It follows that

$$\mathbb{P}[Z^i \neq A_{2t}^i | i \in G] = \mathbb{P}[Z^j \neq A_{2t}^j | j \in G].$$

Let (i_1, \dots, i_k) be representatives of the orbits V/H . Then

$$q(2t) = \max_k \{\mathbb{P}[Z^{i_k} \neq A_{2t}^{i_k} | i_k \in G]\}$$

satisfies the conditions of the claim. \square

3.5.4. Personal estimates of S . Before proceeding to show that S can be well estimated given $\{Z^i\}_{i \in V}$, we note that each agent's opinion is always equal to S with probability at least p . This is true, by definition, at time $t = 0$, and it may not be surprising that this is also the case in later times. However, the proof of this fact is not completely straightforward.

We show in Claim 3.5 above that in the asynchronous case, with probability one there does not exist an infinite sequence of times $t_1 > t_2 > \dots > 0$ such that i updates at time t_i and $i \in N(i+1)$. Let \bar{T} denote an arbitrary choice of clock ring times for which indeed such a chain does not exist. In the synchronous case let \bar{T} be a trivial (probability one) event.

Claim 3.8. *For any \bar{T} , agent i and time t it holds that*

$$\mathbb{P}[A_t^i = S | \bar{T}] \geq p.$$

Proof. Conditioned on \bar{T} , C_t^i is fixed and so A_t^i is a deterministic function of $\{A_0^j : j \in C_t^i\}$. Let $k = |C_t^i|$, and denote this function by $f : \{-1, +1\}^k \rightarrow \{-1, +1\}$.

Clearly, f is monotone, in the sense that

$$f(x_1, \dots, x_{i-1}, x_i, x_{i+1}, \dots, x_k) \leq f(x_1, \dots, x_{i-1}, +1, x_{i+1}, \dots, x_k),$$

and is symmetric in the sense that

$$f(-x_1, -x_2, \dots, -x_k) = -f(x_1, x_2, \dots, x_k).$$

It then follows from Lemma 6.1 in [15] that

$$\mathbb{P}[A_t^i = +1 | \bar{T}, S = +1] \geq p,$$

and so, unconditioned on S , it holds that

$$\mathbb{P}[A_t^i = S | \bar{T}] \geq p. \quad \square$$

3.5.5. Estimating S . As in the setting of Theorem 1.7, let G_0 be a quasi-transitive infinite graph with maximal degree d such that $M_d[G_0, i] < \infty$ for some (\leftrightarrow all) $i \in V_0$. Pick the sets $W_{t,\delta}$ from V_0 , as described above in Section 3.5.2.

Let $H \leq \text{Aut}(G_0)$ act transitively on G_0 , and let $G = (V, E)$ be an infinite connected random subgraph of G_0 with an H -invariant distribution. Let $V_{t,\delta} = W_{t,\delta} \cap V$. Note that $W_{t,\delta}$ and V are independent and are both infinite and H -invariant. Furthermore, as we note above, $W_{t,\delta}$ intersects every orbit $V/\text{Aut}(G_0)$ infinitely. It follows that $V_{t,\delta}$ is almost surely infinite.

Claim 3.9. *Fix an enumeration of the vertices of G_0 , let $\{i_1, i_2, \dots\}$ be the induced enumeration of the vertices in $V_{t,\delta}$, and let*

$$\hat{S}_{t,\delta}^n = \text{sgn} \frac{1}{n} \sum_{k=1}^n A_t^{i_k}.$$

Then

$$\mathbb{P}[\hat{S}_{t,\delta}^n \neq S] \leq e^{-\frac{(p-1/2)^2}{2p}n} + n^2\delta.$$

Proof. Let E be the event that the clock ring times are such that $\{i_1, \dots, i_n\}$ are not causally connected at time t . This happens with probability at least $1 - n^2\delta$. Then conditioned on E , by Claim 3.6, the events ' $A_t^{i_k} = S$ ' are independent.

Now, by Claim 3.8, for all i and t it holds that $\mathbb{P}[A_t^i = S|E] \geq p$. Hence the claim follows by the Chernoff bound. \square

We are now ready to prove Theorem 1.7

Proof of Theorem 1.7. We will prove the theorem by showing that for every $\epsilon > 0$ there exists an \hat{S}_ϵ that is \mathcal{A}_∞ -measurable and such that

$$\mathbb{P}[\hat{S}_\epsilon \neq S] \leq \epsilon.$$

Let N and δ be such that for every t it holds that

$$\mathbb{P}[\hat{S}_{t,\delta}^N \neq S] \leq \epsilon/2, \quad (3.4)$$

where, as in Claim 3.9, $\hat{S}_{t,\delta}^N$ is given by

$$\hat{S}_{t,\delta}^N = \text{sgn} \frac{1}{N} \sum_{k=1}^N A_t^{i_k},$$

with $\{i_1, i_2, \dots\}$ an enumeration of the vertices in $V_{t,\delta}$ which is induced by a fixed enumeration of the vertices of G_0 . The existence of such N and δ satisfying Eq. 3.4 is guaranteed by Claim 3.9.

Let q be a function which satisfies

$$q(2t) \geq \mathbb{P}[Z^i \neq A_{2t}^i], \quad \lim_{t \rightarrow \infty} q(t) = 0,$$

as given in Claim 3.7. Let $T \in 2\mathbb{N}$ be such that $q(T) \leq \epsilon/(2N)$, and define \hat{S}_ϵ by

$$\hat{S}_\epsilon = \text{sgn} \frac{1}{N} \sum_{k=1}^N Z^{i_k},$$

where $\{i_1, i_2, \dots\}$ is the same enumeration used to define $\hat{S}_{T,\delta}^N$. Then

$$\mathbb{P}[\hat{S}_\epsilon \neq \hat{S}_{T,\delta}^N] \leq \mathbb{P}[Z^{i_k} \neq A_T^{i_k} \text{ for some } 1 \leq k \leq N].$$

Now,

$$\mathbb{P}[Z^i \neq A_T^i | i \in V_{t,\delta}] = \mathbb{P}[Z^i \neq A_T^i] \leq q(T).$$

The equality holds since the choice of $V_{t,\delta}$ is independent of the majority dynamics process, and the inequality is simply a reference to the definition of $q(\cdot)$.

Hence by the union bound we have that

$$\mathbb{P}[Z^{i_k} \neq A_T^{i_k} \text{ for some } 1 \leq k \leq N] \leq N \cdot q(T) \leq \epsilon/2,$$

and so

$$\mathbb{P}[\hat{S}_\epsilon \neq \hat{S}_{T,\delta}^N] \leq \epsilon/2.$$

Combining this with Eq. 3.4, we have that

$$\mathbb{P}[\hat{S}_\epsilon \neq S] \leq \epsilon.$$

\square

3.6. Families of uniformly bounded growth graphs. Recall that given $d \geq 3$ and a function $f : \mathbb{N} \rightarrow \mathbb{N}$, $\mathcal{B}(f, d)$ is the family of d -bounded degree graphs G such that $n_r(G, i) \leq f(r)$ for all vertices i in G and $r \in \mathbb{N}$. Recall also that we say that $f : \mathbb{N} \rightarrow \mathbb{N}$ has *slow growth* if $M_d[f] < \infty$.

Fix $d \geq 3$, and let f have slow growth. Let r_0 be the smallest integer for which it holds that

$$\frac{d+1}{d-1} \cdot 2d \cdot \sum_{r>r_0} \left(\frac{d-1}{d+1}\right)^r \cdot f(r) < 1.$$

That is, r_0 is the minimal number that satisfies the condition of Theorem 1.6. Let $B_0(G, i)$ denote the set of vertices at distance at most $r_0 + 2$ from i in G , and note that $|B_0(G, i)|$ is uniformly bounded for all graphs $G \in \mathcal{B}(f, d)$.

Proof of Theorem 1.8. Denote by U_i the event that $A_0^j = S$ for all $j \in B_0(G, i)$. Note that by Theorem 1.6, the event ‘ $Z^i = S$ ’ contains U_i .

Let $p_0 < 1$ be close enough to 1 so that, for any vertex i in any graph $G \in \mathcal{B}(f, d)$, and any $p > p_0$, it holds that

$$\mathbb{P}[U_i] \geq \eta > \frac{1}{2}$$

for some η . This holds, for example, for $p_0 > 1 - 1/(2N)$, where N is a uniform bound on $|B_0(G, i)|$.

Given a graph $G \in \mathcal{B}(f, d)$, let $I_G \subset V$ be a maximal set of vertices such that $B_0(G, i) \cap B_0(G, j) = \emptyset$ for all $i, j \in I_G$. Then the events $\{U_i\}_{i \in I_G}$ are independent, since the events ‘ $A_0^i = S$ ’ are independent.

We first consider the case of finite graphs. Let $\{G_n\}_{n \in \mathbb{N}}$ in $\mathcal{B}(f, d)$ be a sequence of graphs such that $\lim_n |V_n| = \infty$, and let $I_n = I_{G_n}$. Note that $\lim_{n \rightarrow \infty} |I_n| = \infty$, since the diameters of G_n tend to infinity.

Since $\mathbb{P}[U_i] \geq \eta$, and since the events U_i are independent, then by the Chernoff bound we have that

$$\mathbb{P} \left[\sum_{i \in I_n} \mathbb{1}_{U_i} < |I_n|/2 \right] \leq e^{-\frac{(\eta-1/2)^2}{2\eta} |I_n|},$$

and so, since the event ‘ $Z^i = S$ ’ contains U_i , we have that

$$\mathbb{P} \left[\hat{S}_n \neq S \right] \leq e^{-\frac{(\eta-1/2)^2}{2\eta} |I_n|},$$

for

$$\hat{S}_n = \text{sgn} \sum_{i \in I_n} Z^i.$$

Since $\lim_{n \rightarrow \infty} |I_n| = \infty$ it follows that

$$\lim_n \delta(G_n, p) = 0,$$

and we have proved the claim for finite graphs.

Let $G = (V, E) \in \mathcal{B}(f, d)$ be an infinite graph. Then, since the diameter of G is infinite, there exists for every n a set $I_n \subset V$ of size n such that $B_0(G, i) \cap B_0(G, j) = \emptyset$ for all $i, j \in I_n$. By the same reasoning as above, we have that

$$\mathbb{P} \left[\hat{S}_n \neq S \right] \leq e^{-\frac{(\eta-1/2)^2}{2\eta} n},$$

for

$$\hat{S}_n = \operatorname{sgn} \sum_{i \in I_n} Z^i,$$

and so

$$\delta(G, p) = 0.$$

□

3.7. Convergence on a non-slow growth graph. In this section we provide an example of a d -regular graph for which $M_d[G, i] = \infty$ for all $i \in V$, but still the opinion of every agent converges, in the synchronous model, to period at most two, for every set of initial opinions.

Given an odd $d \geq 3$, denote by H the complete bipartite graph $K_{d,d}$ on $2d$ vertices, without a single edge $\{a, b\}$. Let G' be any d -regular graph. Let F be any set of edges whose removal separates G' into finite components. Consider the following process. For an edge $e = \{i, j\} \in F$, remove e from G' , add a copy H_e of H , connect i to a_e (the vertex in H_e which corresponds to a), and connect j to b_e . Repeat this process for all $e \in F$. The resulting graph will be denoted by G .

Lemma 3.10. *For any initial choice of opinions for G , the opinions of all vertices converge, in the synchronous model, to period at most two.*

Since G' was arbitrary, one can easily construct such graphs G with $M_d[G, i] = \infty$.

Proof. We first show that for any copy of H the opinions converge at time at most $t = 2$ to a period at most two. Fix a copy of H , call it also H .

Indeed, if we call one side of H A and the other B , so that $a \in A$, $b \in B$, then at time $t = 1$ all vertices of $A \setminus \{a\}$ have value m_B , the majority of values in B , and similarly vertices in $B \setminus \{b\}$ have value m_A . As $d \geq 3$, both a and b have more neighbors in H than in $G \setminus H$. Thus, at time $t = 2$ all A have opinions m_A , and all B have m_B . Now it is easy to see that for this subgraph we have converged to a period of length at most two.

Since when we remove the copies of H from G we get a union of finite components, after time $t = 2$, the processes in these components are independent and can be thought as applying the dynamics to finite graphs, with some fixed boundary conditions. Thus, the opinions in these components will also converge to a period of length at most two [9].

□

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(O. Tamuz) WEIZMANN INSTITUTE OF SCIENCE, FACULTY OF MATHEMATICS AND COMPUTER SCIENCE, REHOVOT, ISRAEL.

E-mail address, O. Tamuz: omer.tamuz@weizmann.ac.il

(R.J. Tessler) HEBREW UNIVERSITY, EINSTEIN INSTITUTE OF MATHEMATICS, JERUSALEM, ISRAEL.

E-mail address, R.J. Tessler: ran.tessler@mail.huji.ac.il